

CLAIMS

1 1. A method for identifying or disabling at least one traitor receiver with at least one
2 associated unique, compromised decryption key in a broadcast encryption system, comprising:

3 receiving a set of subsets derived from a tree defining leaves, each leaf representing
4 a respective receiver;

5 identifying at least one traitor subset from the set of subsets as containing at least one
6 leaf representing a traitor receiver; and

7 using the traitor subset, identifying or disabling the traitor receiver.

1 2. The method of Claim 1, further comprising:

2 determining whether the traitor subset represents at least one traitor receiver, and if
3 so, dividing the traitor subset into two child sets.

1 3. The method of Claim 2, further comprising determining whether the traitor subset is
2 a member of a frontier set, and if so, removing a complementary subset from the frontier set.

1 4. The method of Claim 1, wherein the act of identifying or disabling includes encoding
2 plural subsets of the set of subsets with a false key.

1 5. The method of Claim 4, further comprising executing a binary search on the set of
2 subsets using probabilities.

1 6. The method of Claim 5, wherein the binary search ends by determining that the
2 difference between a probability p_j of decrypting a message when the first j subsets contain the false
3 key and a probability p_{j-1} of decrypting a message when the first $j-1$ subsets contain the false key is
4 at least equal to a predetermined probability.

1 7. The method of Claim 6, wherein the traitor subset is identified when $|p_{j-1} - p_j| > p/m$,
2 wherein m is the number of subsets in the set of subsets.

1 8. The method of Claim 1, wherein the set of subsets is generated by:
2 assigning each receiver in a group of receivers respective private information I_u ;
3 selecting at least one session encryption key K ;
4 partitioning receivers not in a revoked set R into a set of disjoint subsets S_{11}, \dots, S_{im}
5 having associated subset keys L_{11}, \dots, L_{im} ; and
6 encrypting the session key K and the false key with the subset keys L_{11}, \dots, L_{im} .

1 9. The method of Claim 8, wherein the tree includes a root and plural nodes, each node
2 having an associated key, and wherein each receiver is assigned keys from all nodes in a direct path
3 between a leaf representing the receiver and the root.

1 10. The method of Claim 8, wherein the tree includes a root and plural nodes, each node
2 associated with a set of labels, and wherein each receiver is assigned labels from all nodes hanging
3 from a direct path between the receiver and the root but not from nodes in the direct path.

1 11. The method of Claim 10, wherein the revoked set R defines a spanning tree, and
2 wherein the method includes:

3 initializing a cover tree T as the spanning tree;

4 iteratively removing nodes from the cover tree T and adding nodes to the cover tree
5 T until the cover tree T has at most one node.

1 12. A computer program device, comprising:

2 a computer program storage device including a program of instructions usable by a
3 computer, comprising:

4 logic means for accessing a tree to generate a set of subsets of the tree, the tree
5 including leaves representing at least one traitor device characterized by a compromised key;

6 logic means for encrypting a false key j times and for encrypting a session key m-j
7 times, wherein m is a number of subsets in the set of subsets;

8 logic means responsive to the means for encrypting for identifying a traitor subset; and

9 logic means for using the traitor subset to identify or disable the traitor device.

1 13. The computer program device of Claim 12, further comprising:

2 logic means for determining whether the traitor subset represents at least one traitor
3 device, and if so, dividing the traitor subset into two child sets.

1 14. The computer program device of Claim 13, further comprising logic means for
2 determining whether the traitor subset is a member of a frontier set, and if so, removing a
3 complementary subset from the frontier set.

1 15. The computer program device of Claim 12, further comprising logic means for
2 executing a binary search on the set of subsets using probabilities.

1 16. The computer program device of Claim 15, wherein the binary search ends by
2 determining that the difference between a probability p_j of decrypting a message when the first j
3 subsets contain the false key and a probability p_{j-1} of decrypting a message when the first $j-1$ subsets
4 contain the false key is at least equal to a predetermined probability.

1 17. The computer program device of Claim 16, wherein the traitor subset is identified
2 when $|p_{j-1} - p_j| > p/m$, wherein m is the number of subsets in the set of subsets.

1 18. The method of Claim 12, wherein the set of subsets is generated by logic means
2 including:

3 logic means for assigning each receiver in a group of receivers respective private
4 information I_u ;

5 logic means for selecting at least one session encryption key K;
6 logic means for partitioning receivers not in a revoked set R into a set of disjoint
7 subsets S_{i1}, \dots, S_{im} having associated subset keys L_{i1}, \dots, L_{im} ; and
8 logic means for encrypting the session key K and the false key with the subset keys
9 L_{i1}, \dots, L_{im} .

1 19. The computer program device of Claim 18, wherein the tree includes a root and plural
2 nodes, each node having an associated key, and wherein each receiver is assigned keys from all nodes
3 hanging from a direct path between the receiver and the root but not from nodes in the direct path.

1 20. A computer programmed with instructions to cause the computer to execute method
2 acts including:

3 using a false key to encode plural subsets representing stateless receivers, at least one
4 traitor receiver of which is associated with at least one compromised key that has been
5 obtained by at least one pirate receiver; and

6 using the pirate receiver or a clone thereof, determining the identity of the traitor
7 receiver, or rendering the pirate receiver or clone thereof useless for decrypting data using the
8 compromised key.

1 21. The computer of Claim 20, wherein the subsets define a set of subsets, and the method
2 acts undertaken by the computer further include:

3 receiving the set of subsets derived from a tree defining leaves, each leaf representing
4 a respective receiver;

5 identifying at least one traitor subset from the set of subsets as containing at least one
6 leaf representing the traitor receiver; and

7 using the traitor subset, identifying the traitor receiver.

1 22. The computer of Claim 21, wherein the method acts undertaken by the computer
2 further comprise:

3 determining whether the traitor subset represents at least one traitor receiver, and if
4 so, dividing the traitor subset into two child sets.

1 23. The computer of Claim 22, wherein the method acts undertaken by the computer
2 further comprise determining whether the traitor subset is a member of a frontier set, and if so,
3 removing a complementary subset from the frontier set.

1 24. The computer of Claim 21, wherein the act of identifying includes:
2 encoding plural subsets of the set of subsets with the false key.

1 25. The computer of Claim 24, wherein the method acts undertaken by the computer
2 further comprise executing a binary search on the set of subsets using probabilities.

1 26. The computer of Claim 25, wherein the binary search ends by determining that a
2 probability p_j of decrypting a message when the first j subsets contain the false key is at least equal
3 to a predetermined probability.

1 27. The computer of Claim 26, wherein the traitor subset is identified when $|p_{j-1} - p_j|$
2 $> p/m$, wherein m is the number of subsets in the set of subsets.

1 28. The computer of Claim 21, wherein the set of subsets is generated by:
2 assigning each receiver in a group of receivers respective private information I_u ;
3 selecting at least one session encryption key K ;
4 partitioning receivers not in a revoked set R into a set of disjoint subsets S_{i1}, \dots, S_{im}
5 having associated subset keys L_{i1}, \dots, L_{im} ; and
6 encrypting the session key K and the false key with the subset keys L_{i1}, \dots, L_{im} , wherein
7 the tree includes a root and plural nodes, each node being associated with a set of labels, and
8 wherein each receiver is assigned labels from all nodes hanging from a direct path between
9 the receiver and the root but not from nodes in the direct path.

1 29. The method of Claim 1, further comprising identifying or disabling plural traitor
2 receivers embodied in a clone.

1 30. The method of Claim 1, wherein the act of identifying or disabling includes encoding
2 the first j subsets of the set of subsets with a false key.